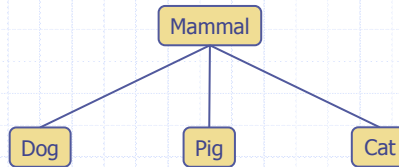


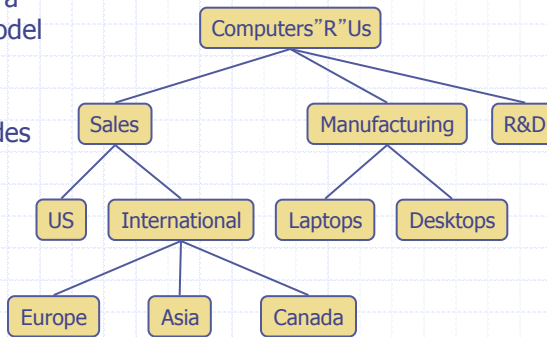
Presentation for use with the textbook **Data Structures and Algorithms in Java, 6th edition**, by M. T. Goodrich, R. Tamassia, and M. H. Goldwasser, Wiley, 2014

Trees



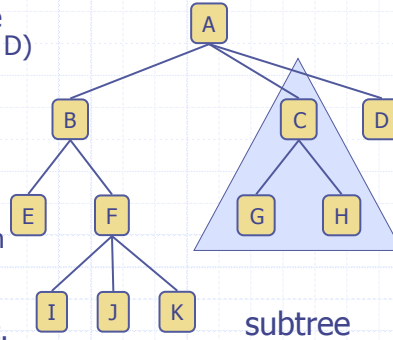
What is a Tree

- In computer science, a tree is an abstract model of a hierarchical structure
- A tree consists of nodes with a parent-child relation
- Applications:
 - Organization charts
 - File systems
 - Programming environments



Tree Terminology

- Root: node without parent (A)
- Internal node: node with at least one child (A, B, C, F)
- External node (a.k.a. leaf): node without children (E, I, J, K, G, H, D)
- Ancestors of a node: parent, grandparent, grand-grandparent, etc.
- Depth of a node: number of ancestors
- Height of a tree: maximum depth of any node (3)
- Descendant of a node: child, grandchild, grand-grandchild, etc.
- Subtree: tree consisting of a node and its descendants



Tree ADT

- We use positions to abstract nodes
- Generic methods:
 - integer `size()`
 - boolean `isEmpty()`
 - Iterator `iterator()`
 - Iterable `positions()`
- Accessor methods:
 - position `root()`
 - position `parent(p)`
 - Iterable `children(p)`
 - Integer `numChildren(p)`
- ◆ Query methods:
 - boolean `isInternal(p)`
 - boolean `isExternal(p)`
 - boolean `isRoot(p)`
- ◆ Additional update methods may be defined by data structures implementing the Tree ADT

Java Interface

Methods for a Tree interface:

```

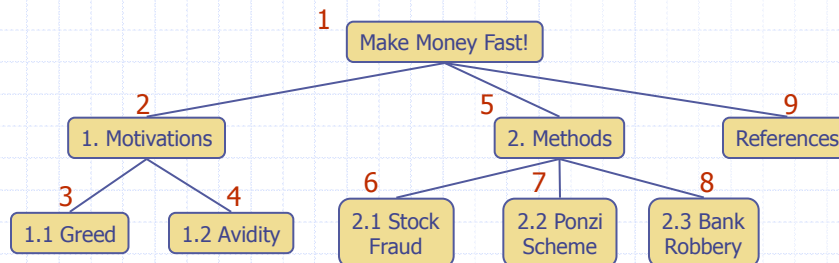
1  /** An interface for a tree where nodes can have an arbitrary number of children. */
2  public interface Tree<E> extends Iterable<E> {
3      Position<E> root();
4      Position<E> parent(Position<E> p) throws IllegalArgumentException;
5      Iterable<Position<E>> children(Position<E> p)
6          throws IllegalArgumentException;
7      int numChildren(Position<E> p) throws IllegalArgumentException;
8      boolean isInternal(Position<E> p) throws IllegalArgumentException;
9      boolean isExternal(Position<E> p) throws IllegalArgumentException;
10     boolean isRoot(Position<E> p) throws IllegalArgumentException;
11     int size();
12     boolean isEmpty();
13     Iterator<E> iterator();
14     Iterable<Position<E>> positions();
15 }

```

Preorder Traversal

- A traversal visits the nodes of a tree in a systematic manner
- In a preorder traversal, a node is visited before its descendants
- Application: print a structured document

Algorithm *preOrder(v)*
visit(v)
for each child *w* of *v*
preorder(w)

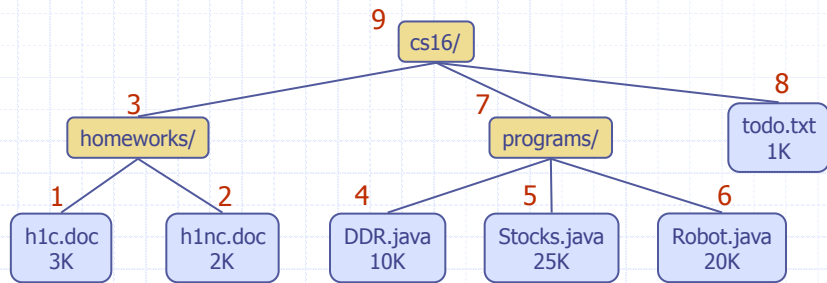


Postorder Traversal

- In a postorder traversal, a node is visited after its descendants
- Application: compute space used by files in a directory and its subdirectories

```

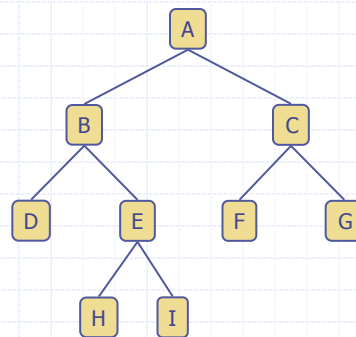
Algorithm postOrder(v)
  for each child w of v
    postOrder(w)
  visit(v)
    
```



Binary Trees

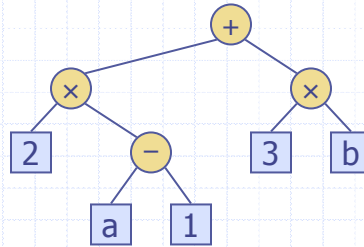
- A binary tree is a tree with the following properties:
 - Each internal node has at most two children (exactly two for **proper** binary trees)
 - The children of a node are an ordered pair
- We call the children of an internal node **left child** and **right child**
- Alternative recursive definition: a binary tree is either
 - a tree consisting of a single node, or
 - a tree whose root has an ordered pair of children, each of which is a binary tree

- Applications:
 - arithmetic expressions
 - decision processes
 - searching



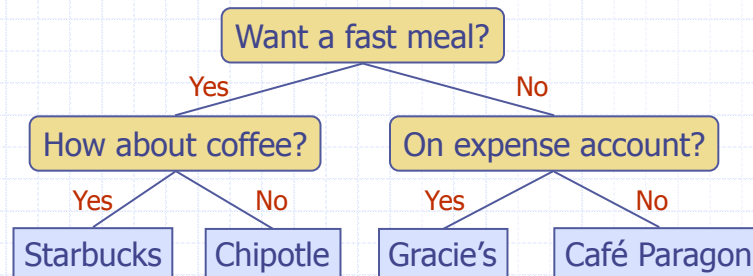
Arithmetic Expression Tree

- Binary tree associated with an arithmetic expression
 - internal nodes: operators
 - external nodes: operands
- Example: arithmetic expression tree for the expression $(2 \times (a - 1) + (3 \times b))$



Decision Tree

- Binary tree associated with a decision process
 - internal nodes: questions with yes/no answer
 - external nodes: decisions
- Example: dining decision



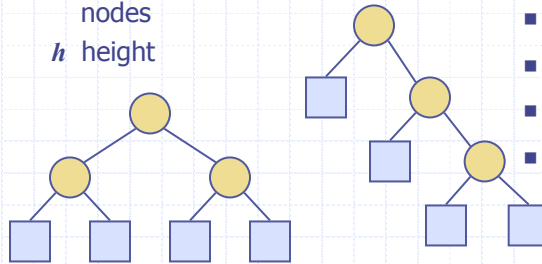
Properties of Proper Binary Trees

□ Notation

- n number of nodes
- e number of external nodes
- i number of internal nodes
- h height

◆ Properties:

- $e = i + 1$
- $n = 2e - 1$
- $h \leq i$
- $h \leq (n - 1)/2$
- $e \leq 2^h$
- $h \geq \log_2 e$
- $h \geq \log_2 (n + 1) - 1$



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BinaryTree ADT

- The **BinaryTree** ADT extends the **Tree** ADT, i.e., it inherits all the methods of the **Tree** ADT
- Additional methods:
 - position **left**(p)
 - position **right**(p)
 - position **sibling**(p)
- The above methods return **null** when there is no left, right, or sibling of p, respectively
- Update methods may be defined by data structures implementing the **BinaryTree** ADT

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Trees

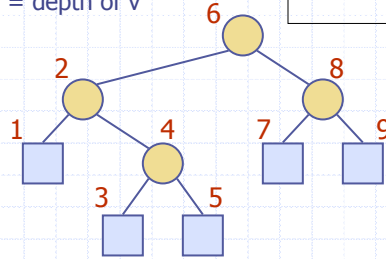
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Inorder Traversal

- In an inorder traversal a node is visited after its left subtree and before its right subtree
- Application: draw a binary tree
 - $x(v)$ = inorder rank of v
 - $y(v)$ = depth of v

```

Algorithm inOrder( $v$ )
  if  $left(v) \neq null$ 
    inOrder( $left(v)$ )
  visit( $v$ )
  if  $right(v) \neq null$ 
    inOrder( $right(v)$ )
    
```



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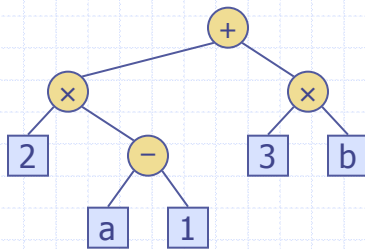
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Print Arithmetic Expressions

- Specialization of an inorder traversal
 - print operand or operator when visiting node
 - print "(" before traversing left subtree
 - print ")" after traversing right subtree

```

Algorithm printExpression( $v$ )
  if  $left(v) \neq null$ 
    print("(' ' )")
    inOrder( $left(v)$ )
  print( $v.element$  ())
  if  $right(v) \neq null$ 
    inOrder( $right(v)$ )
    print("(' ' ' ')")
    
```



$$((2 \times (a - 1)) + (3 \times b))$$

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Trees

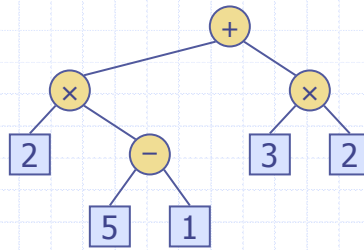
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Evaluate Arithmetic Expressions

- Specialization of a postorder traversal
 - recursive method returning the value of a subtree
 - when visiting an internal node, combine the values of the subtrees

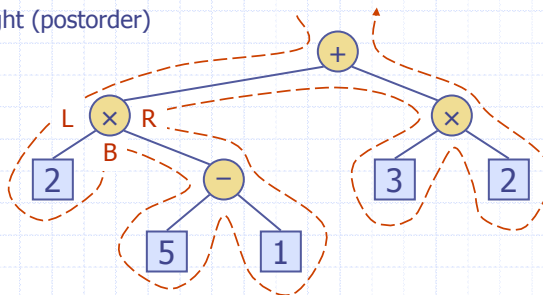
```

Algorithm evalExpr(v)
    if isExternal(v)
        return v.element()
    else
        x ← evalExpr(left(v))
        y ← evalExpr(right(v))
         $\diamond$  ← operator stored at v
        return x  $\diamond$  y
    
```



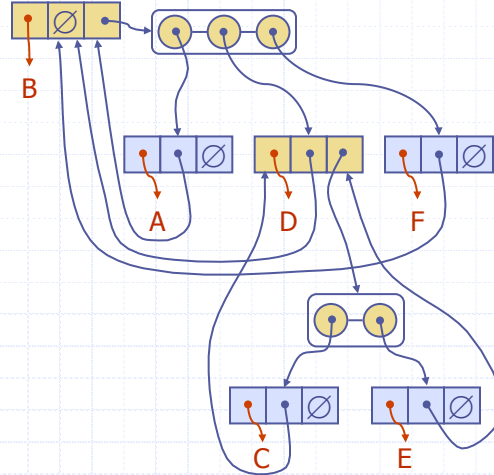
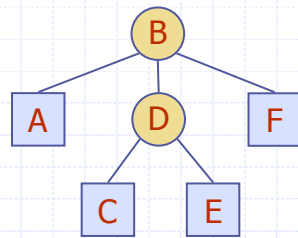
Euler Tour Traversal

- Generic traversal of a binary tree
- Includes a special cases the preorder, postorder and inorder traversals
- Walk around the tree and visit each node three times:
 - on the left (preorder)
 - from below (inorder)
 - on the right (postorder)



Linked Structure for Trees

- A node is represented by an object storing
 - Element
 - Parent node
 - Sequence of children nodes
- Node objects implement the Position ADT



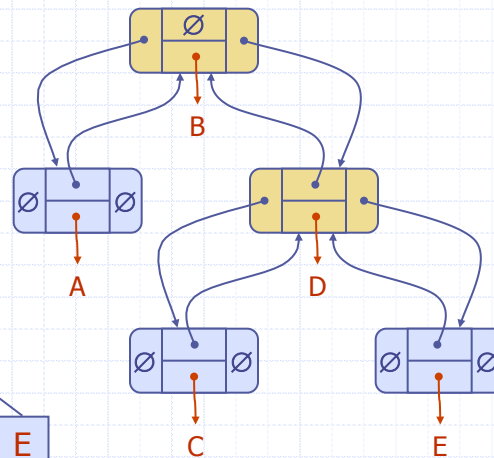
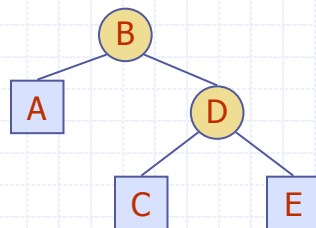
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Trees

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Linked Structure for Binary Trees

- A node is represented by an object storing
 - Element
 - Parent node
 - Left child node
 - Right child node
- Node objects implement the Position ADT



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Trees

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Array-Based Representation of Binary Trees

- Nodes are stored in an array A



Node v is stored at $A[\text{rank}(v)]$

- rank(root) = 0
- if node is the left child of parent(node),
rank(node) = $2 \cdot \text{rank}(\text{parent}(\text{node})) + 1$
- if node is the right child of parent(node),
rank(node) = $2 \cdot \text{rank}(\text{parent}(\text{node})) + 2$

